Math Symbols

| $\wedge$ Conjunction | $\vee$ Disjunction | $\sim$ Negation |
| :--- | :--- | :--- |
| $\rightarrow$ Conditional | $\leftrightarrow$ Biconditional |  |
| $\Rightarrow$ Implication | $\Leftrightarrow$ Equivalence |  |
| $\forall$ For all | $\exists$ There exists | $\in$ in |
| $\therefore$ Therefore | $\square$ QED |  |

Rules of Valid Argumentation Involving Implication Let $P, Q, R$, and $S$ be statements,

| i | $(P \rightarrow Q) \wedge P \Rightarrow Q$ | Modus Ponens (Mode that affirms), used in Direct <br> Proof |
| :---: | :--- | :--- |
| ii | $(P \rightarrow Q) \wedge \sim Q \Rightarrow \sim P$ | Modus Tollens (Mode that denies), used in Proof <br> by Contrapositive |
| iii | $(P \wedge Q) \Rightarrow P$ | Specialization |
| iv | $(P \wedge Q) \Rightarrow Q$ | Specialization |
| v | $P \Rightarrow P \vee Q$ | Addition |
| vii | $(P \vee Q) \wedge \sim P \Rightarrow Q$ | Modus Tollendo Ponens (also called Disjunctive <br> Syllogism. Mode which, by taking away, affirms) |
| ix | $P \leftrightarrow Q \Rightarrow P \rightarrow Q$ | Biconditional-Conditional |
| x | $P \leftrightarrow Q \Rightarrow Q \rightarrow P$ | Biconditional-Conditional |
| xi | $(P \rightarrow Q) \wedge(Q \rightarrow P) \Rightarrow P \leftrightarrow Q$ | Conditional-Biconditional |
| xii | $(P \rightarrow Q) \wedge(Q \rightarrow R) \Rightarrow P \rightarrow R$ | Hypothetical Syllogism |
| xiii | $(P \rightarrow Q) \wedge(R \rightarrow S) \wedge(P \vee R) \Rightarrow(Q \vee S)$ | Constructive Dilema |

Rules of Valid Argumentation Involving Equivalence
Let $P, Q$, and $R$ be statements,

| i | $\sim(\sim P) \Leftrightarrow P$ | Double negation, used in Proof by Contradiction |
| :---: | :--- | :--- |
| ii | $P \vee Q \Leftrightarrow Q \vee P$ | Commutative Law |
| iii | $P \wedge Q \Leftrightarrow Q \wedge P$ | Commutative Law |
| iv | $(P \vee Q) \vee R \Leftrightarrow P \vee(Q \vee R)$ | Associative Law |
| v | $(P \wedge Q) \wedge R \Leftrightarrow P \wedge(Q \wedge R)$ | Associative Law |
| vi | $P \wedge(Q \vee R) \Leftrightarrow(P \wedge Q) \vee(P \wedge R)$ | Distributive Law |
| vii | $P \vee(Q \wedge R) \Leftrightarrow(P \vee Q) \wedge(P \vee R)$ | Distributive Law |
| viii | $P \rightarrow Q \Leftrightarrow \sim P \vee Q$ |  |
| ix | $P \rightarrow Q \Leftrightarrow \sim Q \rightarrow \sim P$ | Contrapositive, used in proofs by Contrapositive |
| x | $P \leftrightarrow Q \Leftrightarrow Q \leftrightarrow P$ |  |
| xi | $P \leftrightarrow Q \Leftrightarrow(P \rightarrow Q) \wedge(Q \rightarrow P)$ |  |
| xii | $\sim(P \wedge Q) \Leftrightarrow \sim P \vee \sim Q$ | De Morgan's Law |
| xiii | $\sim(P \vee Q) \Leftrightarrow \sim P \wedge \sim Q$ | De Morgan's Law |
| xiv | $\sim(P \rightarrow Q) \Leftrightarrow P \wedge \sim Q$ |  |
| xv | $\sim(P \leftrightarrow Q \Leftrightarrow Q) \Leftrightarrow(P \wedge \sim Q) \vee(\sim P \wedge Q)$ |  |

Rules of Valid Argumentation Involving Quantifiers
Let $P(x)$ be a predicate,

| $[(\forall x \in U) P(x)] \Rightarrow P(a)$ where $a$ is arbitrary, | Universal Instantiation |
| :--- | :--- |
| we can chose it to be whatever we want |  |
| $[(\exists x \in U) P(x)] \Rightarrow P(b)$ where $b$ is some particular element of $U$, | Existential Instantiation |
| $b$ can not have appeared before |  |
| $P(c)$ where $c$ is an arbitrary element of $U \Rightarrow[(\forall x \in U) P(x)]$ | Universal Generalization |
| $P(d)$ where $d$ is some particular element of $U \Rightarrow[(\exists x \in U) P(x)]$ | Existential Generalization |

Thanks to Professor McNicholas for this handy one page summary of the Laws of Inference and Equivalence.

